

# Smoothness and Decay Properties of the Limiting Quicksort Density Function

JAMES ALLEN FILL<sup>1</sup>

Department of Mathematical Sciences

The Johns Hopkins University

jimfill@jhu.edu and <http://www.mts.jhu.edu/~fill/>

AND

SVANTE JANSON

Department of Mathematics

Uppsala University

svante.janson@math.uu.se and <http://www.math.uu.se/~svante/>

## ABSTRACT

Using Fourier analysis, we prove that the limiting distribution of the standardized random number of comparisons used by **Quicksort** to sort an array of  $n$  numbers has an everywhere positive and infinitely differentiable density  $f$ , and that each derivative  $f^{(k)}$  enjoys superpolynomial decay at  $\pm\infty$ . In particular, each  $f^{(k)}$  is bounded. Our method is sufficiently computational to prove, for example, that  $f$  is bounded by 16.

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## 1 Introduction and summary

The **Quicksort** algorithm of Hoare [7] is “one of the fastest, the best-known, the most generalized, the most completely analyzed, and the most widely used algorithms for sorting an array of numbers” [2]. **Quicksort** is the standard sorting procedure in **Unix** systems, and Philippe Flajolet, a leader in the field of analysis of algorithms, has noted that it is among “some of the most basic algorithms—the ones that do deserve deep investigation” [4]. Our goal in this introductory section is to review briefly some of what is known about the analysis of **Quicksort** and to summarize how this paper advances that analysis.

The **Quicksort** algorithm for sorting an array of  $n$  numbers is extremely simple to describe. If  $n = 0$  or  $n = 1$ , there is nothing to do. If  $n \geq 2$ , pick a number uniformly at random from the given array. Compare the other numbers to it to partition the remaining numbers into two subarrays. Then recursively invoke **Quicksort** on each of the two subarrays.

Let  $X_n$  denote the (random) number of comparisons required (so that  $X_0 = 0$ ). Then  $X_n$  satisfies the distributional recurrence relation

$$X_n \stackrel{\mathcal{L}}{=} X_{U_n-1} + X_{n-U_n}^* + n - 1, \quad n \geq 1,$$

where  $\stackrel{\mathcal{L}}{=}$  denotes equality in law (i.e., in distribution), and where, on the right,  $U_n$  is distributed uniformly on the set  $\{1, \dots, n\}$ ,  $X_j^* \stackrel{\mathcal{L}}{=} X_j$ , and

$$U_n; X_0, \dots, X_{n-1}; X_0^*, \dots, X_{n-1}^*$$

are all independent.

As is well known and quite easily established, for  $n \geq 0$  we have

$$\mu_n := \mathbf{E}X_n = 2(n+1)H_n - 4n \sim 2n \ln n,$$

where  $H_n := \sum_{k=1}^n k^{-1}$  is the  $n$ th harmonic number and  $\sim$  denotes asymptotic equivalence. It is also routine to compute explicitly the standard deviation of  $X_n$  (see Exercise 6.2.2-8 in [9]), which turns out to be  $\sim n\sqrt{7 - \frac{2}{3}\pi^2}$ .

Consider the standardized variate

$$Y_n := (X_n - \mu_n)/n, \quad n \geq 1.$$

Régnier [11] showed using martingale arguments that  $Y_n \rightarrow Y$  in distribution, with  $Y$  satisfying the distributional identity

$$(1.1) \quad Y \stackrel{\mathcal{L}}{=} UY + (1-U)Z + g(U) =: h_{Y,Z}(U),$$

where

$$(1.2) \quad g(u) := 2u \ln u + 2(1-u) \ln(1-u) + 1,$$

and where, on the right of  $\stackrel{\mathcal{L}}{=}$  in (1.1),  $U$ ,  $Y$ , and  $Z$  are independent, with  $Z \stackrel{\mathcal{L}}{=} Y$  and  $U \sim \text{unif}(0, 1)$ . Rösler [12] showed that (1.1) characterizes the limiting law  $\mathcal{L}(Y)$ , in the precise sense that  $F := \mathcal{L}(Y)$  is the *unique* fixed point of the operator

$$G = \mathcal{L}(V) \mapsto SG := \mathcal{L}(UV + (1 - U)V^* + g(U))$$

(in what should now be obvious notation) subject to

$$\mathbf{E}V = 0, \quad \mathbf{Var}V < \infty.$$

Thus it is clear that fundamental (asymptotic) probabilistic understanding of **Quick-sort**'s behavior relies on fundamental understanding of the limiting distribution  $F$ . In this regard, Rösler [12] showed that

(1.3) the moment generating function (mgf) of  $Y$  is everywhere finite,

and Hennequin [5] [6] and Rösler showed how all the moments of  $Y$  can be pumped out one at a time, though there is no known expression for the mgf nor for the general  $p$ th moment in terms of  $p$ . Tan and Hadjicostas [15] proved that  $F$  has a density  $f$  which is almost everywhere positive, but their proof does not even show whether  $f$  is continuous.

The main goal of this paper is to prove that  $F$  has a density  $f$  which is infinitely differentiable, and that each derivative  $f^{(k)}(y)$  decays as  $y \rightarrow \pm\infty$  more rapidly than any power of  $|y|^{-1}$ : this is our main Theorem 3.1. In particular, it follows that each  $f^{(k)}$  is bounded (cf. Theorem 3.3).

Our main tool will be Fourier analysis. We begin in Section 2 by showing (see Theorem 2.9) that the characteristic function  $\phi$  for  $F$  has rapidly decaying derivatives of every order. Standard arguments reviewed briefly at the outset of Section 3 then immediately carry this result over from  $\phi$  to  $f$ . Finally, in Section 4 we will use the boundedness and continuity of  $f$  to establish an integral equation for  $f$  (Theorem 4.1). As a corollary,  $f$  is everywhere positive (Corollary 4.2).

**Remark 1.1.** (a) Our method is sufficiently computational that we will prove, for example, that  $f$  is bounded by 16. This is not sharp numerically, as Figure 4 of [15] strongly suggests that the maximum value of  $f$  is about  $2/3$ . However, in future work we will rigorously justify (and discuss how to obtain bounds on the error in) the numerical computations used to obtain that figure, and the rather crude bounds on  $f$  and its derivatives obtained in the present paper are needed as a starting point for that more refined work.

(b) Very little is known rigorously about  $f$ . For example, the figure discussed in (a) indicates that  $f$  is unimodal. Can this be proved? Is  $f$  in fact *strongly* unimodal (i.e., log-concave)? What can one say about changes of signs for the derivatives of  $f$ ?

(c) Knessl and Szpankowski [8] purport to prove very sharp estimates of the rates of decay of  $f(y)$  as  $y \rightarrow -\infty$  and as  $y \rightarrow \infty$ . Roughly put, they assert that the left tail of  $f$  decays doubly exponentially (like the tail of an extreme-value density) and that the right tail decays exponentially. But their “results” rely on several unproven assumptions. Among these, for example, is their assumption (59) that

$$\mathbf{E}e^{-\lambda Y} \sim \exp(\alpha\lambda \ln \lambda + \beta\lambda + \gamma \ln \lambda + \delta) \quad \text{as } \lambda \rightarrow \infty$$

for some constants  $\alpha(>0), \beta, \gamma, \delta$ . [Having assumed this, they derive the values of  $\alpha, \gamma$ , and  $\delta$  exactly, and the value of  $\beta$  numerically.]

## 2 Bounds on the limiting Quicksort characteristic function

We will in this section prove the following result on superpolynomial decay of the characteristic function of the limit variable  $Y$ .

**Theorem 2.1.** *For every real  $p \geq 0$  there is a smallest constant  $0 < c_p < \infty$  such that the characteristic function  $\phi(t) \equiv \mathbf{E}e^{itY}$  satisfies*

$$(2.1) \quad |\phi(t)| \leq c_p |t|^{-p} \quad \text{for all } t \in \mathbf{R}.$$

*These best possible constants  $c_p$  satisfy  $c_0 = 1$ ,  $c_{1/2} \leq 2$ ,  $c_{3/4} \leq \sqrt{8\pi}$ ,  $c_1 \leq 4\pi$ ,  $c_{3/2} < 187$ ,  $c_{5/2} < 103215$ ,  $c_{7/2} < 197102280$ , and the relations*

$$(2.2) \quad c_{p_1}^{1/p_1} \leq c_{p_2}^{1/p_2}, \quad 0 < p_1 \leq p_2;$$

$$(2.3) \quad c_{p+1} \leq 2^{p+1} c_p^{1+(1/p)} p/(p-1), \quad p > 1;$$

$$(2.4) \quad c_p \leq 2^{p^2+6p}, \quad p > 0.$$

[The numerical bounds are not sharp (except in the trivial case of  $c_0$ ); they are the best that we can get without too much work, but we expect that substantial improvements are possible.]

*Proof.* The basic approach is to use the fundamental relation (1.1). We will first show, using a method of van der Corput [1, 10], that the characteristic function of  $h_{y,z}(U)$  is bounded by  $2|t|^{-1/2}$  for each  $y, z$ . Mixing, this yields Theorem 2.1 for  $p = 1/2$ . Then we will use another consequence of (1.1), namely, the functional equation

$$(2.5) \quad \phi(t) = \int_{u=0}^1 \phi(ut) \phi((1-u)t) e^{itg(u)} du, \quad t \in \mathbf{R},$$

or rather its consequence

$$(2.6) \quad |\phi(t)| \leq \int_{u=0}^1 |\phi(ut)| |\phi((1-u)t)| du,$$

and obtain successive improvements in the exponent  $p$ .

We give the details as a series of lemmas, beginning with a standard calculus estimate [10]. Note that it suffices to consider  $t > 0$  in the proofs because  $\phi(-t) = \overline{\phi(t)}$  and thus  $|\phi(-t)| = |\phi(t)|$ . Note also that the best constants satisfy  $c_p = \sup_{t>0} t^p |\phi(t)|$  (although we do not know in advance of proving Theorem 2.1 that these are finite), and thus  $c_p^{1/p} = \sup_{t>0} t |\phi(t)|^{1/p}$ , which clearly satisfies (2.2) because  $|\phi(t)| \leq 1$ .

**Lemma 2.2.** *Suppose that a function  $h$  is twice continuously differentiable on an open interval  $(a, b)$  with*

$$h'(x) \geq c > 0 \quad \text{and} \quad h''(x) \geq 0 \quad \text{for } x \in (a, b).$$

Then

$$\left| \int_{x=a}^b e^{ith(x)} dx \right| \leq \frac{2}{ct} \text{ for all } t > 0.$$

*Proof.* By considering subintervals  $(a + \varepsilon, b - \varepsilon)$  and letting  $\varepsilon \rightarrow 0$ , we may without loss of generality assume that  $h$  is defined and twice differentiable at the endpoints, too. Then, using integration by parts, we calculate

$$\begin{aligned} \int_{x=a}^b e^{ith(x)} dx &= \frac{1}{it} \int_{x=a}^b \left[ \frac{d}{dx} e^{ith(x)} \right] \frac{dx}{h'(x)} \\ &= \frac{1}{it} \left\{ \frac{e^{ith(x)}}{h'(x)} \Big|_{x=a}^b - \int_{x=a}^b e^{ith(x)} d \left( \frac{1}{h'(x)} \right) \right\}. \end{aligned}$$

So

$$\begin{aligned} \left| \int_{x=a}^b e^{ith(x)} dx \right| &\leq \frac{1}{t} \left\{ \left( \frac{1}{h'(b)} + \frac{1}{h'(a)} \right) + \int_{x=a}^b \left| d \left( \frac{1}{h'(x)} \right) \right| dx \right\} \\ &= \frac{1}{t} \left\{ \left( \frac{1}{h'(b)} + \frac{1}{h'(a)} \right) + \int_{x=a}^b \left[ -d \left( \frac{1}{h'(x)} \right) \right] dx \right\} \\ &= \frac{1}{t} \left\{ \left( \frac{1}{h'(b)} + \frac{1}{h'(a)} \right) + \left( \frac{1}{h'(a)} - \frac{1}{h'(b)} \right) \right\} \\ &= \frac{2}{th'(a)} \leq \frac{2}{ct}. \end{aligned} \quad \square$$

**Lemma 2.3.** For any real numbers  $y$  and  $z$ , the random variable  $h_{y,z}(U)$  defined by (1.1) satisfies

$$|\mathbf{E} e^{ith_{y,z}(U)}| \leq 2|t|^{-1/2}.$$

*Proof.* We will apply Lemma 2.2, taking  $h$  to be  $h_{y,z}$ . Observe that

$$h''_{y,z}(u) = 2 \left( \frac{1}{u} + \frac{1}{1-u} \right) = \frac{2}{u(1-u)} \geq 8 \text{ for } u \in (0, 1)$$

and that

$$h'_{y,z}(u) = 0 \text{ if and only if } u = \alpha_{y,z} := \frac{1}{1 + \exp(\frac{1}{2}(y-z))} \in (0, 1).$$

Let  $t > 0$  and  $\gamma > 0$ . If in Lemma 2.2 we take  $a := \alpha_{y,z} + \gamma t^{-1/2}$  and  $b := 1$ , and assume that  $a < b$ , then note

$$h'(u) = h'_{y,z}(u) = \int_{x=\alpha_{y,z}}^u h''_{y,z}(x) dx \geq 8(u - \alpha_{y,z}) \geq 8\gamma t^{-1/2} \text{ for all } u \in (a, b).$$

So, by Lemma 2.2,

$$\left| \int_{u=\alpha_{y,z} + \gamma t^{-1/2}}^1 e^{ith_{y,z}(u)} du \right| \leq \frac{2}{t} [8\gamma t^{-1/2}]^{-1} = \frac{1}{4\gamma} t^{-1/2}.$$

Trivially,

$$\left| \int_{u=\alpha_{y,z}}^{\alpha_{y,z}+\gamma t^{-1/2}} e^{ith_{y,z}(u)} du \right| \leq \gamma t^{-1/2},$$

so we can conclude

$$\left| \int_{u=\alpha_{y,z}}^1 e^{ith_{y,z}(u)} du \right| \leq [(4\gamma)^{-1} + \gamma] t^{-1/2}.$$

This result is trivially also true when  $a = \alpha_{y,z} + \gamma t^{-1/2} \geq b = 1$ , so it holds for all  $t, \gamma > 0$ . The optimal choice of  $\gamma$  here is  $\gamma_{\text{opt}} = 1/2$ , which yields

$$\left| \int_{u=\alpha_{y,z}}^1 e^{ith_{y,z}(u)} du \right| \leq t^{-1/2} \quad \text{for all } t > 0.$$

Similarly, for example by considering  $u \mapsto h(1-u)$ ,

$$\left| \int_0^{\alpha_{y,z}} e^{ith_{y,z}(u)} du \right| \leq t^{-1/2} \quad \text{for all } t > 0,$$

and we conclude that the lemma holds for all  $t > 0$ , and thus for all real  $t$ .  $\square$

**Lemma 2.4.** *For any real  $t$ ,  $|\phi(t)| \leq 2|t|^{-1/2}$ .*

*Proof.* Lemma 2.3 shows that

$$\left| \mathbf{E} \left( e^{ith_{Y,Z}(U)} \mid Y, Z \right) \right| \leq 2|t|^{-1/2}$$

and thus

$$|\phi(t)| = \left| \mathbf{E} e^{ith_{Y,Z}(U)} \right| \leq \mathbf{E} \left| \mathbf{E} \left( e^{ith_{Y,Z}(U)} \mid Y, Z \right) \right| \leq 2|t|^{-1/2}. \quad \square$$

The preceding lemma is the case  $p = 1/2$  of Theorem 2.1. We now improve the exponent.

**Lemma 2.5.** *Let  $0 < p < 1$ . Then*

$$c_{2p} \leq \frac{[\Gamma(1-p)]^2}{\Gamma(2-2p)} c_p^2.$$

*Proof.* By (2.6) and the definition of  $c_p$ ,

$$|\phi(t)| \leq \int_{u=0}^1 c_p^2 |ut|^{-p} |(1-u)t|^{-p} du = c_p^2 |t|^{-2p} \int_{u=0}^1 u^{-p} (1-u)^{-p} du,$$

and the result follows by evaluating the beta integral.  $\square$

In particular, recalling  $\Gamma(1/2) = \sqrt{\pi}$ , Lemmas 2.4 and 2.5 yield

$$(2.7) \quad |\phi(t)| \leq \frac{4\pi}{|t|}.$$

This proves (2.1) for  $p = 1$ , with  $c_1 \leq 4\pi$ , and thus by (2.2) for every  $p \leq 1$  with  $c_p \leq (4\pi)^p$ ; applying Lemma 2.5 again, we obtain the finiteness of  $c_p$  in (2.1) for all  $p < 2$ . Somewhat better numerical bounds are obtained for  $1/2 < p < 1$  by taking a geometric average between the cases  $p = 1/2$  and  $p = 1$ : the inequality

$$|\phi(t)| \leq (2t^{-1/2})^{2-2p}(4\pi t^{-1})^{2p-1} = 2^{2p}\pi^{2p-1}t^{-p}, \quad t > 0,$$

shows that  $c_p \leq 2^{2p}\pi^{2p-1}$ ,  $1/2 \leq p \leq 1$ . In particular, we have  $c_{3/4} \leq \sqrt{8\pi}$ , and thus, by Lemma 2.5,  $c_{3/2} \leq 8\pi^{1/2}[\Gamma(1/4)]^2 < 186.4 < 187$ .

**Lemma 2.6.** *Let  $p > 1$ . Then*

$$c_{p+1} \leq 2^{p+1}c_p^{1+(1/p)}p/(p-1).$$

*Proof.* Assume that  $t \geq 2c_p^{1/p}$ . Then, again using (2.6),

$$\begin{aligned} |\phi(t)| &\leq \int_{u=0}^1 \min\left(\frac{c_p}{(ut)^p}, 1\right) \min\left(\frac{c_p}{[(1-u)t]^p}, 1\right) du \\ &= 2 \int_{u=0}^{c_p^{1/p}t^{-1}} \frac{c_p}{[(1-u)t]^p} du + \int_{u=c_p^{1/p}t^{-1}}^{1-c_p^{1/p}t^{-1}} \frac{c_p^2}{[u(1-u)t^2]^p} du \\ &\leq \frac{2}{[1-c_p^{1/p}t^{-1}]^p} \frac{c_p^{1+(1/p)}}{t^{p+1}} + 2 \frac{c_p^2}{t^{2p}} \int_{u=c_p^{1/p}t^{-1}}^{1/2} \frac{du}{[u(1-u)]^p} \\ &\leq \frac{2}{(1/2)^p} c_p^{1+(1/p)} t^{-(p+1)} + \frac{2}{(1/2)^p} \frac{c_p^2}{t^{2p}} \int_{u=c_p^{1/p}t^{-1}}^{1/2} u^{-p} du \\ &\leq 2^{p+1} \left\{ c_p^{1+(1/p)} t^{-(p+1)} + \frac{1}{p-1} c_p^2 t^{-2p} [c_p^{1/p}t^{-1}]^{-(p-1)} \right\} \\ &= 2^{p+1} c_p^{1+(1/p)} \frac{p}{p-1} t^{-(p+1)}. \end{aligned}$$

We have derived the desired bound for all  $t \geq 2c_p^{1/p}$ . But also, for all  $0 < t < 2c_p^{1/p}$ , we have

$$2^{p+1} c_p^{1+(1/p)} \frac{p}{p-1} t^{-(p+1)} \geq \frac{p}{p-1} \geq 1 \geq |\phi(t)|,$$

so the estimate holds for all  $t > 0$ . □

Lemma 2.6 completes the proof of finiteness of every  $c_p$  in (2.1) (by induction), and of the estimate (2.3). The bound for  $c_{3/2}$  obtained above now shows (using `Maple`) that  $c_{5/2} < 103215$ , which then gives  $c_{7/2} < 197102280$ .

We can rewrite (2.3) as

$$\begin{aligned} c_{p+1}^{1/(p+1)} &\leq 2c_p^{1/p} \left(1 + \frac{1}{p-1}\right)^{1/(p+1)} \leq 2c_p^{1/p} \exp\left(\frac{1}{(p-1)(p+1)}\right) \\ &= 2c_p^{1/p} \exp\left(\frac{1}{2(p-1)} - \frac{1}{2(p+1)}\right). \end{aligned}$$

Hence, by induction, if  $p = n + \frac{5}{2}$  for a nonnegative integer  $n$ , then

$$c_p^{1/p} \leq 2^n c_{5/2}^{2/5} e^{(1/3)+(1/5)} = C 2^n,$$

where  $C := 2^{-5/2} e^{8/15} c_{5/2}^{2/5} < 30.6 < 2^5$ , using the above estimate of  $c_{5/2}$ . Consequently,  $c_p^{1/p} < 2^{p+5}$  when  $p = n + \frac{5}{2}$ . For general  $p > 3/2$  we now use (2.2) with  $p_1 = p$  and  $p_2 = \lceil p - \frac{5}{2} \rceil + \frac{5}{2}$ , obtaining  $c_p^{1/p} < 2^{p_2+5} < 2^{p+6}$ ; the case  $p \leq 3/2$  follows from (2.2) and the estimate  $c_{3/2}^{2/3} < 33 < 2^6$ . This completes the proof of (2.4) and hence of Theorem 2.1.  $\square$

**Remark 2.7.** We used (1.1) in two different ways. In the first step we conditioned on the values of  $Y$  and  $Z$ , while in the inductive steps we conditioned on  $U$ .

**Remark 2.8.** A variety of other bounds are possible. For example, if we begin with the inequality (2.7), use (2.6), and proceed just as in the proof of Lemma 2.6, we can easily derive the following result in the case  $t \geq 8\pi$ :

$$(2.8) \quad |\phi(t)| \leq \frac{32\pi^2}{t^2} \left( \ln\left(\frac{t}{4\pi}\right) + 2 \right) \leq \frac{32\pi^2 \ln t}{t^2} \quad \text{for all } t \geq 1.72.$$

The result is trivial for  $1.72 \leq t < 8\pi$ , since then the bounds exceed unity.

Since  $Y$  has finite moments of all orders [recall (1.3)], the characteristic function  $\phi$  is infinitely differentiable. Theorem 2.1 implies a rapid decrease of all derivatives, too.

**Theorem 2.9.** *For each real  $p \geq 0$  and integer  $k \geq 0$ , there is a constant  $c_{p,k}$  such that*

$$|\phi^{(k)}(t)| \leq c_{p,k} |t|^{-p} \quad \text{for all } t \in \mathbf{R}.$$

*Proof.* The case  $k = 0$  is Theorem 2.1, and the case  $p = 0$  follows by  $|\phi^{(k)}(t)| \leq \mathbf{E}|Y|^k$ . The remaining cases follow from these cases by induction on  $k$  and the following calculus lemma.  $\square$

**Lemma 2.10.** *Suppose that  $g$  is a complex-valued function on  $(0, \infty)$  and that  $A, B, p > 0$  are such that  $|g(t)| \leq At^{-p}$  and  $|g''(t)| \leq B$  for all  $t > 0$ . Then  $|g'(t)| \leq 2\sqrt{AB}t^{-p/2}$ .*



*Proof.* Fix  $t > 0$  and let  $\theta = \arg(g'(t))$ . For  $s > t$ ,

$$\operatorname{Re}(e^{-i\theta}g'(s)) \geq \operatorname{Re}(e^{-i\theta}g'(t)) - |g'(s) - g'(t)| \geq |g'(t)| - B(s - t)$$

and thus, integrating from  $t$  to  $t_1 := t + (|g'(t)|/B)$ ,

$$\begin{aligned} \operatorname{Re}(e^{-i\theta}(g(t_1) - g(t))) &\geq \int_t^{t_1} (|g'(t)| - B(s - t)) ds \\ &= (t_1 - t)|g'(t)| - \frac{1}{2}B(t_1 - t)^2 = |g'(t)|^2/(2B). \end{aligned}$$

Consequently,

$$|g'(t)|^2/(2B) \leq |g(t)| + |g(t_1)| \leq 2At^{-p},$$

and the result follows.  $\square$

In other words, the characteristic function  $\phi$  belongs to the class  $\mathcal{S}$  of infinitely differentiable functions that, together with all derivatives, decrease more rapidly than any power. (This is the important class of test functions for tempered distributions, introduced by Schwartz [14]; it is often called the class of *rapidly decreasing*  $C^\infty$  functions.)

### 3 The limiting Quicksort density $f$ and its derivatives

We can now improve the result by Tan and Hadjicostas [15] on existence of a density  $f$  for  $Y$ . It is an immediate consequence of Theorem 2.1, with  $p = 0$  and  $p = 2$ , say, that the characteristic function  $\phi$  is integrable over the real line. It is well-known—see, e.g., [3, Theorem XV.3.3]—that this implies that  $Y$  has a bounded continuous density  $f$  given by the Fourier inversion formula

$$(3.1) \quad f(x) = \frac{1}{2\pi} \int_{t=-\infty}^{\infty} e^{-itx} \phi(t) dt, \quad x \in \mathbf{R}.$$

Moreover, using Theorem 2.1 with  $p = k + 2$ , we see that  $t^k \phi(t)$  is also integrable for each integer  $k \geq 0$ , which by a standard argument (cf. [3, Section XV.4]) shows that  $f$  is infinitely smooth, with a  $k$ th derivative ( $k \geq 0$ ) given by

$$(3.2) \quad f^{(k)}(x) = \frac{1}{2\pi} \int_{t=-\infty}^{\infty} (-it)^k e^{-itx} \phi(t) dt, \quad x \in \mathbf{R}.$$

It follows further that the derivatives are bounded, with

$$(3.3) \quad \sup_x |f^{(k)}(x)| \leq \frac{1}{2\pi} \int_{t=-\infty}^{\infty} |t|^k |\phi(t)| dt \quad (k \geq 0),$$

and these bounds in turn can be estimated using Theorem 2.1. Moreover, as is well known [14], [13, Theorem 7.4], an extension of this argument shows that the class  $\mathcal{S}$  discussed at the end of Section 2 is preserved by the Fourier transform, and thus Theorem 2.9 implies that  $f \in \mathcal{S}$ :

**Theorem 3.1.** *The Quicksort limiting distribution has an infinitely differentiable density function  $f$ . For each real  $p \geq 0$  and integer  $k \geq 0$ , there is a constant  $C_{p,k}$  such that*

$$|f^{(k)}(x)| \leq C_{p,k}|x|^{-p} \quad \text{for all } x \in \mathbf{R}. \quad \square$$

For numerical bounds on  $f$ , we can use (3.3) with  $k = 0$  and Theorem 2.1 for several different  $p$  (in different intervals); for example, using  $p = 0, 1/2, 1, 3/2, 5/2, 7/2$ , and taking  $t_1 = 4, t_2 = 4\pi^2, t_3 = (187/(4\pi))^2, t_4 = 103215/187, t_5 = 197102280/103215$ ,

$$\begin{aligned} f(x) &\leq \frac{1}{2\pi} \int_{t=-\infty}^{\infty} |\phi(t)| dt = \frac{1}{\pi} \int_{t=0}^{\infty} |\phi(t)| dt \\ &\leq \frac{1}{\pi} \int_{t=0}^{\infty} \min(1, 2t^{-1/2}, 4\pi t^{-1}, 187t^{-3/2}, 103215t^{-5/2}, 197102280t^{-7/2}) dt \\ (3.4) \quad &= \frac{1}{\pi} \left( \int_{t=0}^{t_1} dt + \int_{t=t_1}^{t_2} 2t^{-1/2} dt + \int_{t=t_2}^{t_3} 4\pi t^{-1} dt + \int_{t=t_3}^{t_4} 187t^{-3/2} dt \right. \\ &\quad \left. + \int_{t=t_4}^{t_5} 103215t^{-5/2} dt + \int_{t=t_5}^{\infty} 197102280t^{-7/2} dt \right) \\ &\leq 18.2. \end{aligned}$$

**Remark 3.2.** We can do somewhat better by using the first bound in (2.8) over the interval  $(103.18, 1984)$  instead of (as above) Theorem 2.1 with  $p = 1, 3/2, 5/2, 7/2$  over  $(103.18, t_3), (t_3, t_4), (t_4, t_5), (t_5, 1984)$ , respectively. This gives

$$f(x) < 15.3.$$

Similarly, (3.3) with  $k = 1$  and the same estimates of  $|\phi(t)|$  as in (3.4) yield

$$|f'(x)| \leq \frac{1}{2\pi} \int_{t=-\infty}^{\infty} |t| |\phi(t)| dt = \frac{1}{\pi} \int_{t=0}^{\infty} t |\phi(t)| dt < 3652.1,$$

which can be reduced to 2492.1 by proceeding as in Remark 3.2. The bound can be further improved to 2465.9 by using also  $p = 9/2$ .

Somewhat better bounds are obtained by using more values of  $p$  in the estimates of the integrals, but the improvements obtained in this way seem to be slight. We summarize the bounds we have obtained.

**Theorem 3.3.** *The limiting Quicksort density function  $f$  satisfies  $\max_x f(x) < 16$  and  $\max_x |f'(x)| < 2466$ .*  $\square$

The numerical bounds obtained here are far from sharp; examination of Figure 4 of [15] suggests that  $\max f < 1$  and  $\max |f'| < 2$ . Our present technique cannot hope to produce a better bound on  $f$  than  $4/\pi > 1.27$ , since neither Lemma 2.3 nor (2.6) can improve on the bound  $|\phi(t)| \leq 1$  for  $|t| \leq 4$ . Further, *no* technique based on (3.3) can hope to do better than the actual value of  $(2\pi)^{-1} \int_{t=-\infty}^{\infty} |\phi(t)| dt$ , which from cursory examination of Figure 6 of [15] appears to be about 2.

## 4 An integral equation for the density $f$

Our estimates are readily used to justify rigorously the following functional equation.

**Theorem 4.1.** *The continuous limiting Quicksort density  $f$  satisfies (pointwise) the integral equation*

$$f(x) = \int_{u=0}^1 \int_{y \in \mathbf{R}} f(y) f\left(\frac{x - g(u) - (1-u)y}{u}\right) \frac{1}{u} dy du, \quad x \in \mathbf{R},$$

where  $g(\cdot)$  is as in (1.2).

*Proof.* For each  $u$  with  $0 < u < 1$ , the random variable

$$(4.1) \quad uY + (1-u)Z + g(u)$$

[with notation as in (1.1)] has the density function

$$(4.2) \quad f_u(x) := \int_{z \in \mathbf{R}} f(z) f\left(\frac{x - g(u) - (1-u)z}{u}\right) \frac{1}{u} dz,$$

where the integral converges for each  $x$  since, using Theorem 3.3, the integrand is bounded by  $f(z)(\max f)/u \leq 16f(z)/u$ ; dominated convergence using the continuity of  $f$  and the same bound shows further that  $f_u$  is continuous.

This argument yields the bound  $f_u(x) \leq 16/u$ , and since  $f_u = f_{1-u}$  by symmetry in (4.1), we have  $f_u(x) \leq 16/\max(u, 1-u) \leq 32$ . This uniform bound, (1.1), and dominated convergence again imply that  $\int_0^1 f_u(x) du$  is a continuous density for  $Y$ , and thus equals  $f(x)$  for every  $x$ .  $\square$

It was shown in [15] that  $f$  is positive almost everywhere; we now can improve this by removing the qualifier “almost.”

**Corollary 4.2.** *The continuous limiting Quicksort density function is everywhere positive.*

*Proof.* We again use the notation (4.2) from the proof of Theorem 4.1. Fix  $x \in \mathbf{R}$  and  $u \in (0, 1)$ . Since  $f$  is almost everywhere positive [15], the integrand in (4.2) is positive almost everywhere. Therefore  $f_u(x) > 0$ . Now we integrate over  $u \in (0, 1)$  to conclude that  $f(x) > 0$ .  $\square$

Alternatively, Corollary 4.2 can be derived directly from Theorem 4.1, without recourse to [15]. Indeed, if  $f(y_0) > 0$  and  $u_0 \in (0, 1)$ , set  $x = y_0 + g(y_0)$ ; then the integrand in the double integral for  $f(x)$  in Theorem 4.1 is positive for  $(u, y)$  equal to  $(u_0, y_0)$ , and therefore, by continuity, also in some small neighborhood thereof. It follows that  $f(y_0 + g(y_0)) > 0$ . Since  $u_0$  is arbitrary and the image of  $(0, 1)$  under  $g$  is  $(-(2 \ln 2 - 1), 1)$ , an open interval containing the origin, Corollary 4.2 follows readily.

**Remark 4.3.** In future work, we will use arguments similar to those of this paper, together with other arguments, to show that when one applies the method of successive substitutions to the integral equation in Theorem 4.1, the iterates enjoy exponential-rate uniform convergence to  $f$ . This will settle an issue raised in the third paragraph of Section 3 in [15].

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